

# Priority Queues and Heaps

Metodický koncept k efektivní podpoře klíčových odborných kompetencí s využitím cizího jazyka ATCZ62 - CLIL jako výuková strategie na vysoké škole

**Interreg**   
EVROPSKÁ UNIE  
**Rakousko-Česká republika**  
Evropský fond pro regionální rozvoj



**Europäische Union**  
**Evropská unie**  
Europäischer Fonds für  
regionale Entwicklung  
Evropský fond pro  
regionální rozvoj



**UNIVERSITY**  
**OF APPLIED SCIENCES**  
**UPPER AUSTRIA**

# Priority Queue – ADT

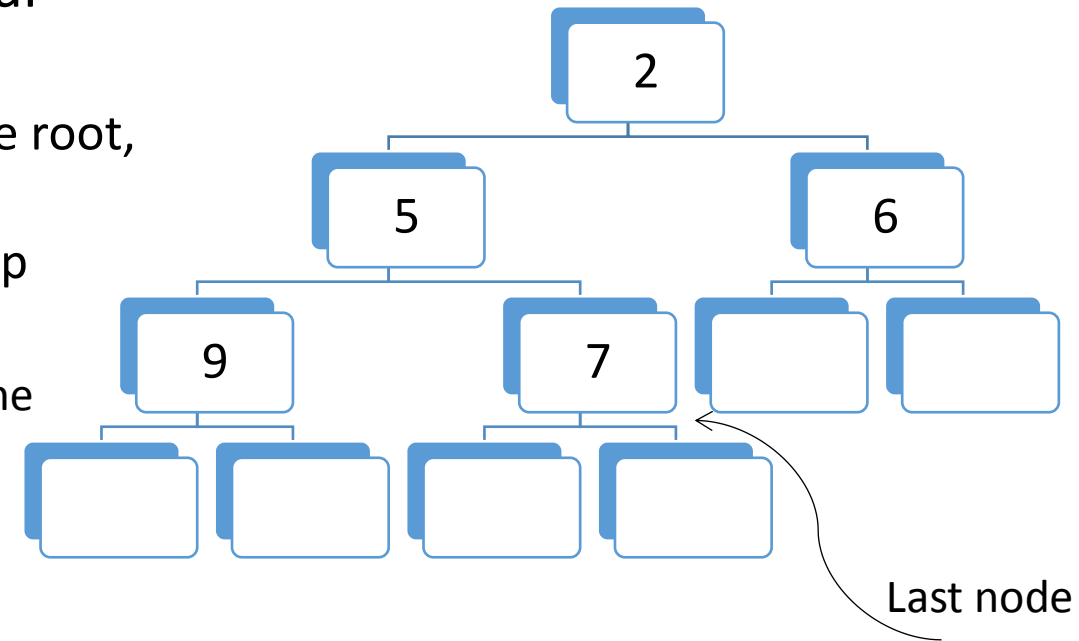
- A priority queue stores a collection of items
- An item is a pair (key, element)
- Main methods of the Priority Queue ADT
  - **insertItem(k, o)** - inserts an item with key k and element o
  - **removeMin()** - removes the item with smallest key and returns its element
- Additional methods
  - **minKey(k, o)**      **minElement()**      **size()**      **isEmpty()**
- Applications:
  - Standby flyers, Auctions, Stock market

# Priority Queue

- Keys in a priority queue can be arbitrary objects on which an order is defined
- Two distinct items in a priority queue can have the same key
- Mathematical concept of total order relation  $\leq$ 
  - Reflexive property:  $x \leq x$
  - Antisymmetric property:  $x \leq y \wedge y \leq x \Rightarrow x = y$
  - Transitive property:  $x \leq y \wedge y \leq z \Rightarrow x \leq z$
- Comparator – ADT
  - A comparator encapsulates the action of comparing two objects according to a given total order relation

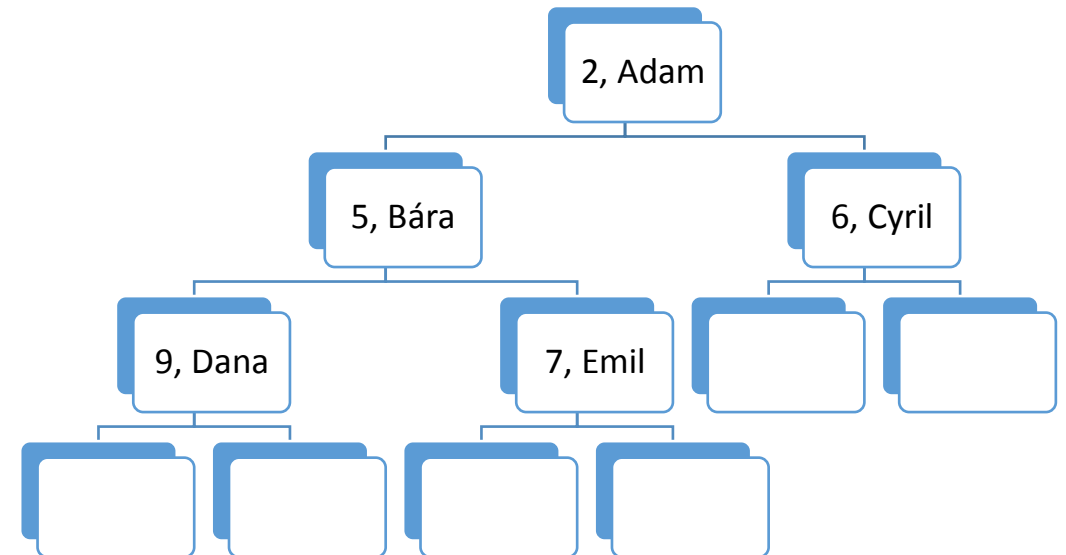
# Heap

- A heap is a binary tree storing keys at its internal nodes and satisfying the following properties:
  - Heap-Order: for every internal node  $v$  other than the root,  $key(v) \geq key(parent(v))$
  - Complete Binary Tree: let  $h$  be the height of the heap
    - for  $i = 0, \dots, h - 1$ , there are  $2^i$  nodes of depth  $i$
    - at depth  $h - 1$ , the internal nodes are to the left of the external nodes
- The last node of a heap is the rightmost internal node of depth  $h - 1$



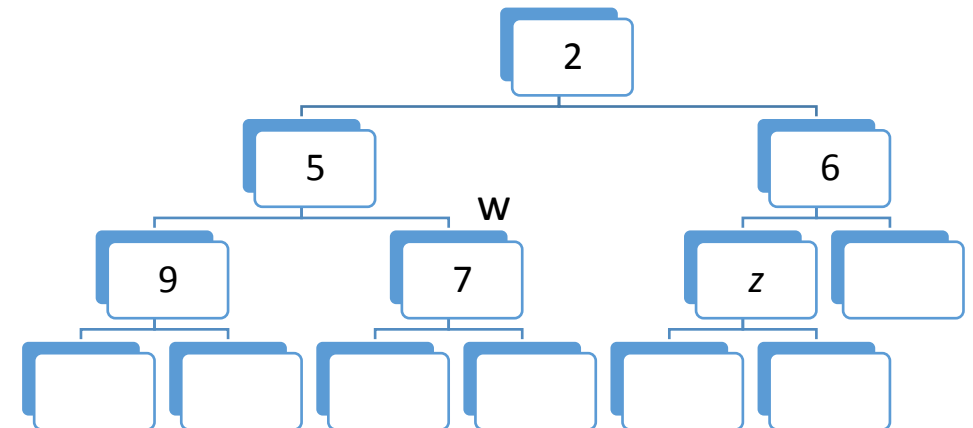
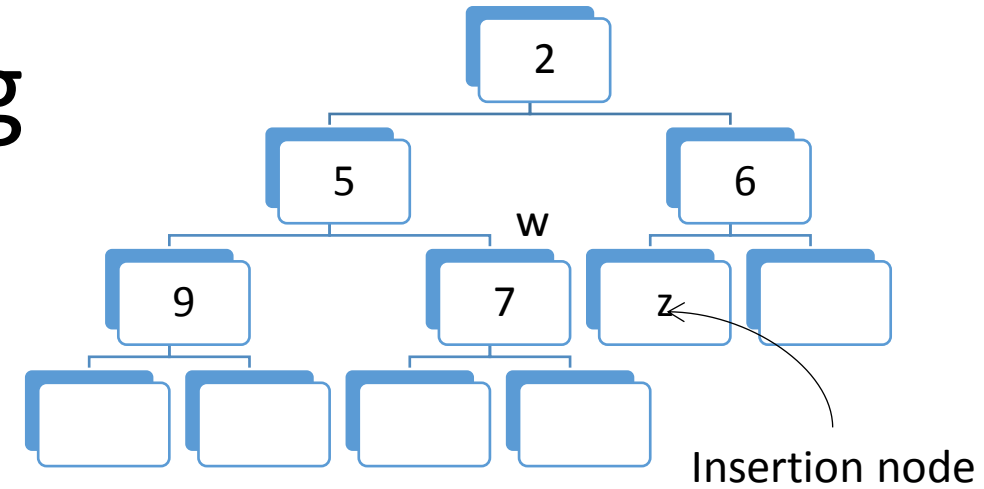
# Heaps and Priority Queue

- We can use a heap to implement a priority queue
- We store a (key, element) item at each internal node
- We keep track of the position of the last node



# Heap – inserting and deleting

- The insertion algorithm consists of three steps
  - Find the insertion node  $z$  (the new last node)
  - Store  $k$  at  $z$  and expand  $z$  into an internal node
  - Restore the heap-order property (discussed next)
- The removal algorithm consists of three steps
  - Replace the root key with the key of the last node  $w$
  - Compress  $w$  and its children into a leaf
  - Restore the heap-order property (discussed next)



# Restoring the heap order

- **upheap()**

- After the insertion of a new key  $k$ , the heap-order property may be violated
- Algorithm upheap restores the heap-order property by swapping  $k$  along an upward path from the insertion node
- Upheap terminates when the key  $k$  reaches the root or a node whose parent has a key smaller than or equal to  $k$

- **downheap()**

- After replacing the root key with the key  $k$  of the last node, the heap-order property may be violated
- Algorithm downheap restores the heap-order property by swapping key  $k$  along a downward path from the root
- Upheap terminates when key  $k$  reaches a leaf or a node whose children have keys greater than or equal to  $k$