

Graph theory

Metodický koncept k efektivní podpoře klíčových odborných kompetencí s využitím cizího jazyka ATCZ62 - CLIL jako výuková strategie na vysoké škole

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Graph theory

- A graph is a pair (V, E) , where
 - V is a set of nodes, called vertices
 - E is a collection of pairs of vertices, called edges
 - Vertices and edges are positions and store elements
- Types of edges
 - Directed – ordered pair of vertices (u, v) , first vertex u is the origin, second vertex v is the destination
 - Undirected - unordered pair of vertices (u, v)
 - Loops – edge that connects a vertex to itself
 - Multiple edges – between vertices (u, v) is more than one edge

Graph theory

- Types of graphs
 - Directed – all edges are direct
 - Undirected – all edges are undirected
 - Multigraph – contains multiple edges
- Terminology
 - End vertices (or endpoints) of an edge
 - Edges incident on a vertex
 - Adjacent vertices
 - Degree of a vertex
 - Parallel edges
 - Self-loop

Graph theory

- Path
 - sequence of alternating vertices and edges
 - begins with a vertex
 - ends with a vertex
 - each edge is preceded and followed by its endpoints
- Simple path
 - path such that all its vertices and edges are distinct
- Cycle
 - circular sequence of alternating vertices and edges
 - each edge is preceded and followed by its endpoints
- Simple cycle
 - cycle such that all its vertices and edges are distinct

Graph theory

- Electronic circuits
 - Printed circuit board
 - Integrated circuit
- Transportation networks
 - Highway network
 - Flight network
- Computer networks
 - Local area network
 - Internet
 - Web
- Databases
 - Entity-relationship diagram

Graph – ADT

- Accessor methods
 - **aVertex()**
 - **incidentEdges(v)**
 - **endVertices(e)**
 - **isDirected(e)**
 - **origin(e)**
 - **destination(e)**
 - **opposite(v,e)**
 - **areAdjacent(v,w)**
- Update methods
 - **insertVertex(o)**
 - **insertEdge(v, w, o)**
 - **insertDirectedEdge(v, w, o)**
 - **removeVertex(v)**
 - **removeEdge(e)**
- Generic methods
 - **numVertices()**
 - **numEdges()**
 - **vertices()**
 - **edges()**

Graph – DFS – depth-first search

- for traversing or searching tree or graph data structures. One starts at the root (selecting some arbitrary node as the root in the case of a graph) and explores as far as possible along each branch before backtracking.

procedure DFS-iterative(G, v):

 let S be a stack

$S.push(v)$

 while S is not empty

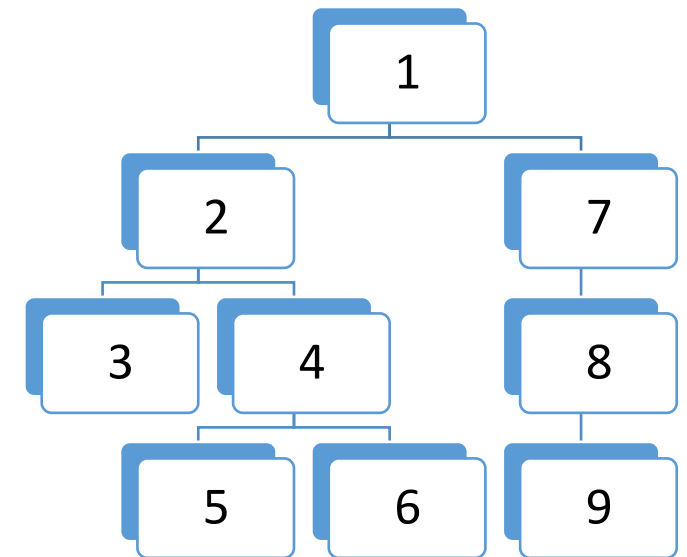
$v = S.pop()$

 if v is not labeled as discovered:

 label v as discovered

 for all edges from v to w in $G.adjacentEdges(v)$ do

$S.push(w)$



Graph – BFS – Breadth-first search

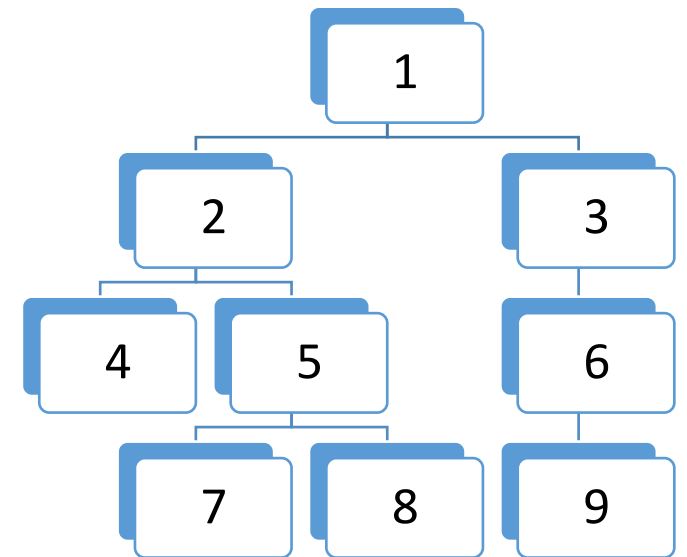
- traversing or searching tree or graph data structures. It starts at the tree root (or some arbitrary node of a graph, sometimes referred to as a 'search key') and explores the neighbor nodes first, before moving to the next level neighbors.

Breadth-First-Search(Graph, root):

```
create empty set S
create empty queue Q
add root to S
Q.enqueue(root)
while Q is not empty:
  current = Q.dequeue()
  if current is the goal:
    return current
```

for each node n that is adjacent to current:

```
if n is not in S:
  add n to S
  n.parent = current
  Q.enqueue(n)
```



Graph – shortest path

- Dijkstra's algorithm
 - non-negative weights on the edges
 - $O(|V|^2 + |E|)$ – V number of vertices, E number of edges
- Bellman-Ford algorithm
 - Graph can have negative edges
 - $O(V \cdot E)$ – slower than Dijkstra's alg.
- Floyd-Warshall algorithm
 - Directed graph with non-negative edges
 - Find shortest path between all vertices
 - Time complexity – $O(V^3)$, memory complexity - $O(V^2)$

Graph – shortest path

- Johnson's algorithm
 - find the shortest paths between all pairs of vertices in a sparse, edge weighted, directed graph. It allows some of the edge weights to be negative numbers
 - $O(V^2 \log_2(V) + VE)$